

1. Multiplying two 2-bit unsigned numbers (H.+H. 8.14, p.493)

| | | | | | | |
|-----------------|-----------|-----------|------|------|------|------|
| Mult. table: | | $c_1 c_0$ | 00 | 01 | 11 | 10 |
| | $d_1 d_0$ | 00 | 0000 | 0000 | 0000 | 0000 |
| | 01 | 0000 | 0001 | 0011 | 0010 | |
| | 11 | 0000 | 0011 | 1001 | 0110 | |
| | 10 | 0000 | 0010 | 0110 | 0100 | |

$= b_3 b_2 b_1 b_0$

Karnaugh maps for each bit :

1-bit b_0

| | | | | | |
|-----------|-----------|----|----|----|----|
| | $c_1 c_0$ | 00 | 01 | 11 | 10 |
| $d_1 d_0$ | 00 | 0 | 0 | 0 | 0 |
| 01 | 0 | 1 | 1 | 0 | |
| 11 | 0 | 1 | 1 | 0 | |
| 10 | 0 | 0 | 0 | 0 | |

$b_0 = c_0 \cdot d_0$



2-bit b_1

| | | | | | |
|-----------|-----------|----|----|----|----|
| | $c_1 c_0$ | 00 | 01 | 11 | 10 |
| $d_1 d_0$ | 00 | 0 | 0 | 0 | 0 |
| 01 | 0 | 0 | 1 | 1 | |
| 11 | 0 | 1 | 0 | 1 | |
| 10 | 0 | 1 | 1 | 0 | |

$b_1 = c_1 d_0 \oplus d_1 c_0$



4-bit b_2

| | | | | | |
|-----------|-----------|----|----|----|----|
| | $c_1 c_0$ | 00 | 01 | 11 | 10 |
| $d_1 d_0$ | 00 | 0 | 0 | 0 | 0 |
| 01 | 0 | 0 | 0 | 0 | |
| 11 | 0 | 0 | 0 | 1 | |
| 10 | 0 | 0 | 1 | 1 | |

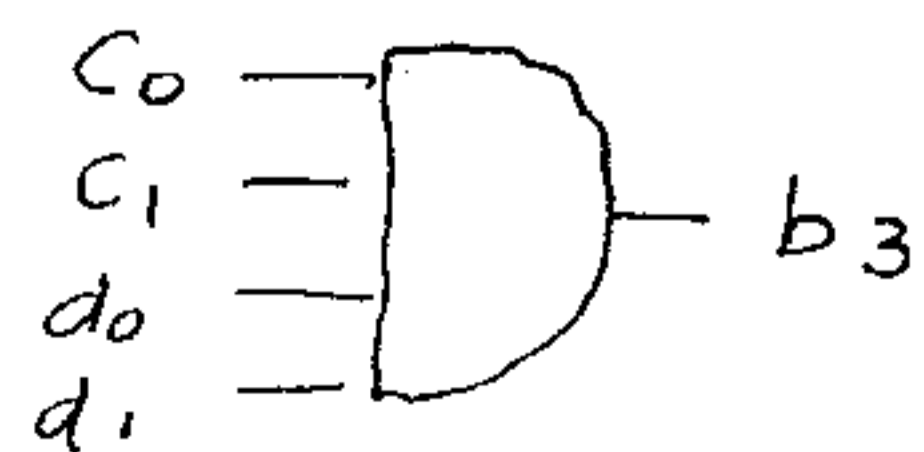
$b_2 = c_1 d_1 \overline{c_0 d_0}$



8-bit b_3

| | | | | | |
|-----------|-----------|----|----|----|----|
| | $c_1 c_0$ | 00 | 01 | 11 | 10 |
| $d_1 d_0$ | 00 | 0 | 0 | 0 | 0 |
| 01 | 0 | 0 | 0 | 0 | |
| 11 | 0 | 0 | 1 | 0 | |
| 10 | 0 | 0 | 0 | 0 | |

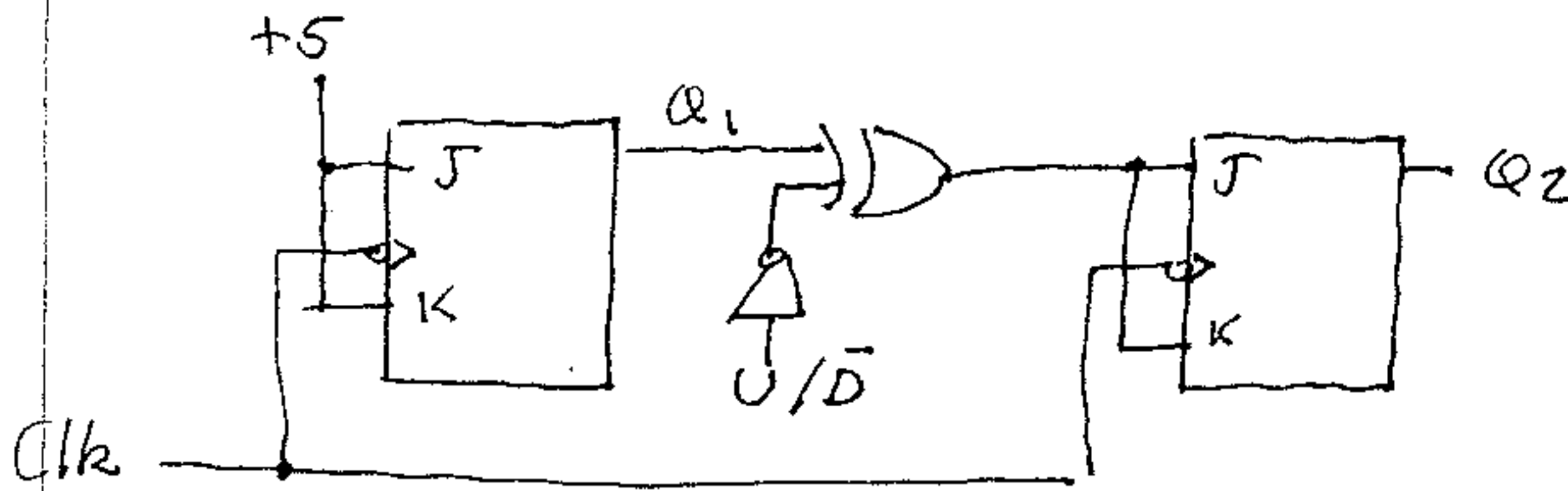
$b_3 = c_1 c_0 d_1 d_0$



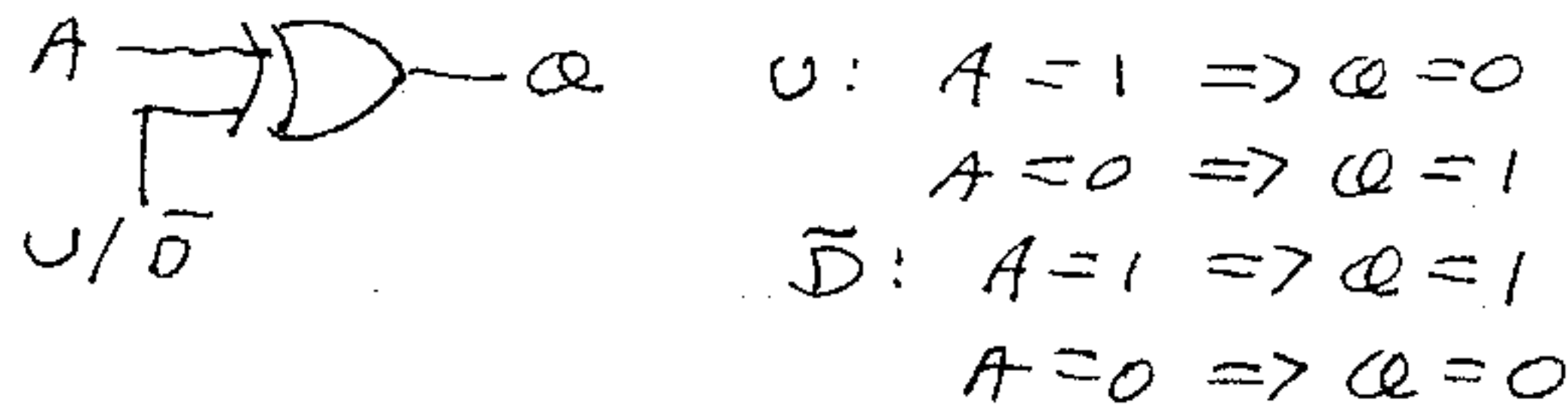
② Review Exercise # 2

6/4/09

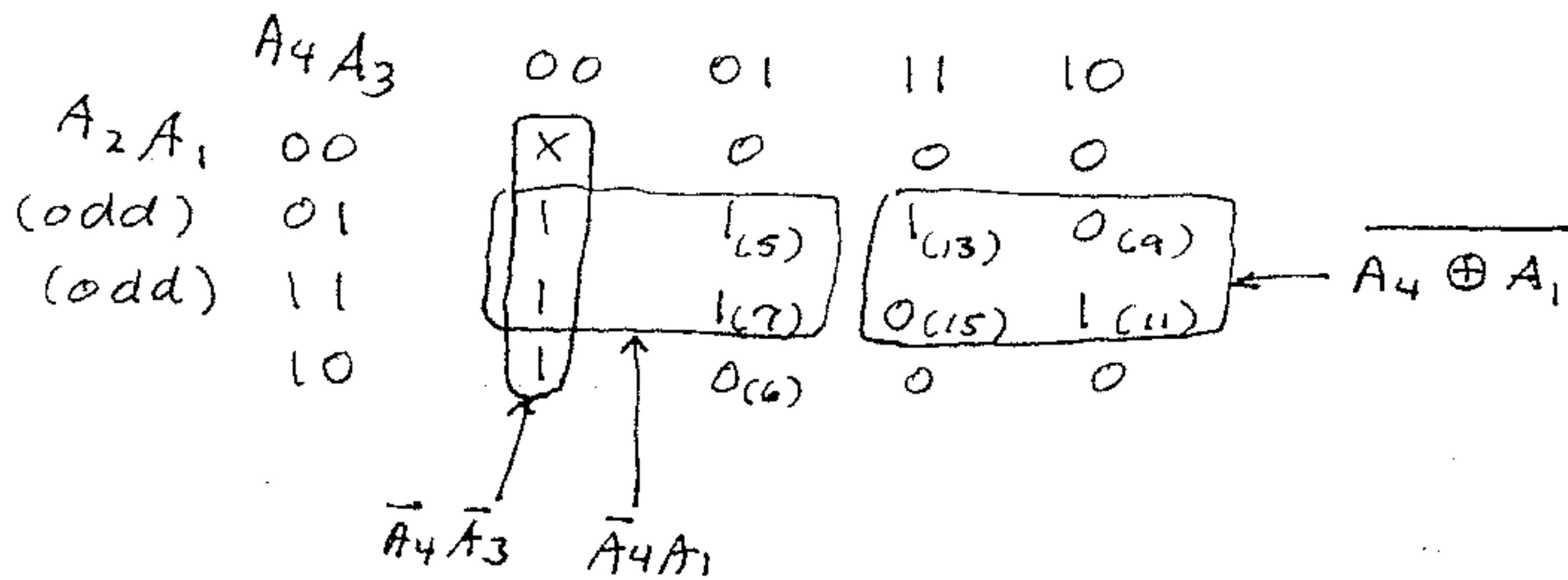
2. Synchronous 2-bit U/D counter (H.+H., 8.25, p.514):



U: counts up
 \bar{D} : counts down

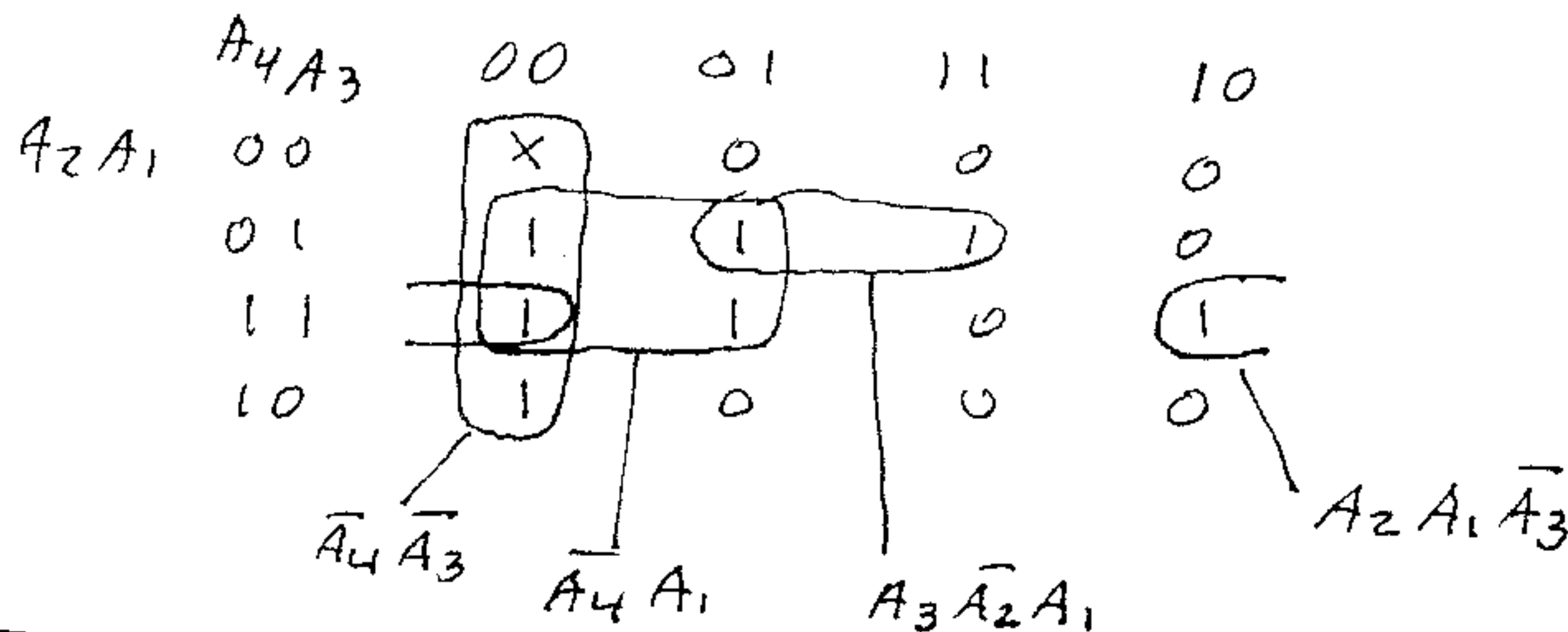


3. Prime number finder (1 through 15):



$$P = \bar{A}_4 (A_1 + \bar{A}_3) + \overline{A_4 \oplus A_1}$$

This is one of several solutions. Another is



So
$$P = \bar{A}_4 (\bar{A}_3 + A_1) + A_1 (\bar{A}_3 A_2 + \bar{A}_2 A_3) = \bar{A}_4 (A_3 + A_1) + A_1 (A_2 \oplus A_3)$$